

## Design Considerations for Packet Networks supporting Synchronous Ethernet and IEEE 1588v2

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#### **Some Quotes for Consideration**

"A synchronous transmission layer will always be there, everywhere"

"Only the transmission layer can deal with synchronization, it cannot be done elsewhere"

- Transmission Engineer

"Why do you need sync??? This is legacy!
Sync is going the way of the dinosaurs"

"When everything has moved to IP and Ethernet, you don't need sync anymore"

- Packet Network Engineer



## **Food for Thought**

Can we reconcile both circuit and packet worlds and create something "better"?

Synchronization can bring a lot to the packet world

- Circuit emulation and radio hand-off support
- Very accurate ToD (network monitoring, transaction management, etc)

Networks are changing – some decisions to be made:

- Do we need sync everywhere ?
- Is sync part of the infrastructure?
- Is sync a service?
  - Or both ?



## **Agenda**

- 1 Packet-based Architectures
- 2 Deploying SyncE
- 3 Deploying IEEE1588v2
- 4 Deploying SyncE and IEEE1588v2 together



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# Packet-based Architectures Change is the Only Constant

#### Packet networks have more diverse topologies than SDH

- Mesh, rings, dual homed rings, stars, dual homed stars, chains, ...
- Number of nodes per island (ex.: ring) ?
- Access vs Aggregation vs Core

#### What packet technologies are used

- Ethernet, MPLS, IP, but also flavors (ETHoSDH, MPLS-TP)
- And associated protocols

#### Who owns what

- Mobile carrier, wireline carrier, wholesale
- Trusted and un-trusted zones
- Requires boundaries with adequate features

## What is (to be) deployed

- Existing (recent) network
- Greenfield deployment
- Network extension

**Every network is different** 



# Packet-based Architectures Sync Strategy

#### Synchronization strategy options

- No network assistance
  - End-to-end transparent 1588v2, spot insertion of frequency/phase at the edge
- Partial network assistance
  - Boundary and transparent clocking support on selected nodes
- Full network assistance
  - Layer-1 transport of Frequency with SyncE
  - IEEE1588v2 with TC (and BC)
  - Including hybrid mode

#### Definition of "network assistance"

- Clock transfer/maintenance capabilities
- QoS enforcement
- Traffic Engineering
- Monitoring



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#### **Deploying SyncE**

#### **Challenge #1: Network Design**

## Challenge:

ALL network components in the path need to support SyncE

## Reality check: Network Boundaries

- SDH to SyncE boundaries still not well proven (clock quality, traceability, loops)
- What about other boundaries (IEEE1588v2)
- Boundaries need expensive oscillators
- Where to start SyncE (in the core?)
- Complexity of meshed networks (not ring-only any more, like SDH)

## SyncE will be deployed in an ever-changing environment

Different from SDH - can be deployed link by link (careful design & validation)

This is layer 1= hardware upgrades

Virtualization ("I want my own clock")



#### **Deploying SyncE**

## **Challenge #2: OAM**

Existing SSUs to support ESMC (traceability, loops, ...)

#### Boundaries

- SDH/SyncE
- OTN/SyncE

#### **ESMC** Performance

- When under 100% load at 100 GE
- Convergence in a distributed chassis

#### Adding TLVs will

- Be beneficial for monitoring
- Increase complexity (performance?)
- Introduce interop issues between SyncE vendors and with SSUs
- Phase support will increase this even further (see next slide)



#### **Deploying SyncE**

## **Challenge #3: Phase Support using ESMC**

#### ESMC re-use is a good target, however

- Why build another protocol for phase support ?
- What does really matter: the protocol, or the end-result?
- How different from IEEE1588 with BC/TC ?
  - Could be less flexible (more embedded into systems)
- Still expensive oscillator needed at every hop (or every line card)

#### Potential issues

- Phase accuracy on a distributed chassis (to be designed into the PFE)
- Inter-working with IEEE1588 (TC, or BC for non-Eth boundaries)
- Will we need a hardware upgrade again ?

#### Caveats

- Less flexible than IEEE1588 (non Ethernet portions of the network ?)
- Virtualization ("I want my own clock")

## SyncE is one part of the puzzle



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Challenge: Network Design in an "un-deterministic" packet world

#### Ultimate target

- Phase accuracy (reach the microsecond level)
- Increase hop count between Master and Slaves
  - while maintaining frequency / phase accuracy

#### Jitter and wander

- PDV: QoS
- Network load: Traffic Engineering & CAC

#### Experience being built\*

- Microwave
- xDSL asymmetric bandwidth
- Network design for phase support

\* not detailed here for the sake of time



Challenge #1: Engineer for Jitter Part 1: Latency

PTP: 64-128 B

reference: 128B

#### Similarity

L2 vs L3

IPv4 vs IPv6 vs MPLS LER/LSR

IP vs L2/L3 VPNs

Centralized vs distributed

No QoS: Max goes to the roof

QoS or RFC2544: similar

What matters is the baseline (min/avg)

 The more information, the better (max to be limited)

#### Fully loaded system

- All ports
- Large switching / routing table
- Impact of other processing

Latency (microseconds)	Min/Avg	Max
1 GE - NPU	200	300
1 GE – L2 (centralized)	5	10
1 GE – L3 (centralized)	5	10
1 GE – L3 (distributed)	10-15	20
10 GE – L2 (centralized)	2	4
10 GE – L3 (centralized)	2	4
10 GE – L3 (distributed)	15-20	30

Vendor Dependent!



Challenge #1: Engineer for Jitter Part 2: QoS choices

#### Strict Priority Queue

- Always served (CAC needed) no extra delay
- PTP + CES (+ VoIP) (+OAM) packet sizes small and consistent
- Some max of SPQ depends on the rest of queuing

#### High Priority Queue

- Credit could be negative: could be starved / buffered
- PTP: high drop probability, low buffer size
- VoIP, Business High, OAM

#### **Network Control Queue**

- Routing protocols have different requirements (could be buffered)
- Different profile from PTP traffic

N packets on the driver after the queue: MTU has an impact (less on 10GE) – check jumbo frames



Vendor

Dependent!

Challenge #1: Engineer for Jitter Part 3: Implementations

#### Demystifying:

- "Switching matrix" distributed systems are more sophisticated
- ETHoSDH may not be so perfect

#### To be checked

Distributed systems (between line cards)



#### QoS with SPQ can compensate for bursts

- Avoid peaks of temporary buffering
- Avoid packet drops

#### What also matters is the long term PDV

Engineer traffic along with QoS (see next topic)



Challenge #1: Engineer for Jitter Part 4: QoS and TC

Good quality NEs down to the access are a good thing anyway

Needed for other things (Business High, VoIP, LTE, etc)

#### Is TC a way to compensate

- For poor quality switches
- For the use of QoS
- Is that really simpler

#### Balance complexity / cost of

- Hardware-based QoS devices with SPQ
- TC capable devices

#### TC is the ultimate goal

- Will eventually be needed for high accuracy phase
- Not yet clear if needed everywhere
- Centralized vs distributed systems
  - requires intelligent calculation/distribution of residence time across Packet Forwarding Engines
- TC at 10GE / 100GE



**Challenge #2: Engineer for Wander** Part 1: Symmetry

#### Forward / reverse path

- Layer 3 routing protocols may be asymmetric, but MPLS solves the problem
- No more issues with asymmetric delays due to data rate steps (1GE, 10GE)
- Potential inter-work with Network Monitoring tools (TWAMP) to calibrate it

Check the NEs have the same behavior in upstream / downstream (QoS impact)

#### Physical networks:

- use fiber instead of copper
- xDSL: we have to live with it slaves to cope with it

What matters is that the asymmetry keeps stable



Challenge #2: Engineer for Wander Part 2: Network Load

#### New rules of NGMNs

- Network convergence
- Multiservice

#### Link load

Core: 50%

Edge/Access: US not at 100%

Fiber P2P star: DS 100%

Ring or MW chain: 100% on first DS link, then less and less

L2 or L3 protocols: little control (shortest path at best)

#### MPLS: brings control

- MPLS DiffServ Traffic Engineering to control congestion over paths
- MPLS CAC to control traffic classes
- Ramp up can better be managed

#### Wholesale islands

Require strong SLAs



Challenge #2: Engineer for Wander Part 3: Delay jumps

Network failures, re-routing, operational mistakes

- There are also some planned/known delay jumps
- Special case of AMP

What matters: a delay jump is stable and deterministic

MPLS DiffServ Traffic Engineering to manage delay jumps

- TE path computation
- Detection
  - MPLS OAM P2P and E2E
- Protection
  - Fast Reroute (FRR): node link path
  - Make before break, Non-revert mode



# Deploying IEEE1588v2 The role of MPLS for Sync – Also Multicast

#### Usual assumptions:

- Too complex to deploy and manage
- Access networks are not multicast capable

Converged networks require multicast anyway (multiplay)

#### Multicast MPLS is

- Simpler to manage
- More deterministic than IP multicast

#### Traffic Engineering can be used for Multicast also

Including protection

#### Benefits

- Scale the PTP Master (direct impact on cost and stability)
- Increase PTP performance



# Deploying IEEE1588v2 The role of MPLS for Sync – Sync as a Service

What is Sync: in a packet network, is sync part of transmission, or a service, or part of network control? Something else?

Usual assumption: MPLS cannot be everywhere, too expensive

#### New rules apply

- MPLS moving from routing to transmission
  - "The Purple Line" Kireeti Kompella
- MPLS is already in the core
- MPLS is already in the aggregation
- MPLS can be as far as in the access (MPLS-TP is an example)

#### MPLS to bring flexibility

- Decouple transport and service
  - « Seamless MPLS » (draft-leymann-mpls-seamless-mpls-00.txt; Oct. 20th, 2009)
  - Sync as a service
  - 1588 Masters are Service Nodes
- Direct and protected connections to the Masters
  - E-Line, L2 VPN, L3 VPN
- Network engineering without constraints: place / move your Service Nodes freely



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## **Comparing SyncE and IEEE1588v2**

	Synchronous Ethernet	IEEE 1588v2
Ease of Deployment	Disruptive	Non-Disruptive (except if TC and BC are needed)
Cost of Deployment	Inexpensive for greenfield Expensive for upgrade (Capex +Opex)	Inexpensive (except if TC and BC are needed)
Ubiquity	No (Ethernet only)	Yes (Ethernet, xDSL, Microwave, Cable, FTTx, femtocell, etc)
Frequency/Phase Future-proof	No - Frequency only (Except if ESMC used for Phase But HW upgrades)	Yes (Frequency & Phase)
Deterministic behavior	Yes	Under certain conditions (better if using BC and TC)
Requires Careful Engineering	Yes	Yes
Virtualization	No	Yes (if overlay)



## **Combining SyncE and IEEE1588v2**

#### Why combined SyncE & IEEE1588v2?

Algorithms to enhance quality of both; clock failover; cross boundaries

#### It is not only a slave problem

- Also in any boundary node (islands)
- It makes sense to plan BC and SyncE in the same NE
- Will eventually make sense to combine TC and SyncE in the same NE

#### May well be required in every Node / POP

Where applications are to be supported

#### Open question: what is the best way to implement?

- Combined algorithms ("hybrid mode")
- Combined protocols (ESMC & TC ?), including OAM
- Hardware upgrade?

## SyncE & IEEE1588v2 are complementary





## **Key Takeaways**

Adding clocking to packet networks can be gradual

Actual deployment will take time

MPLS and Traffic Engineering can play an important role

e.g. "Sync as a Service"





